

Optimal arrangements of n points on a sphere and in a circle

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IMFM/TCS

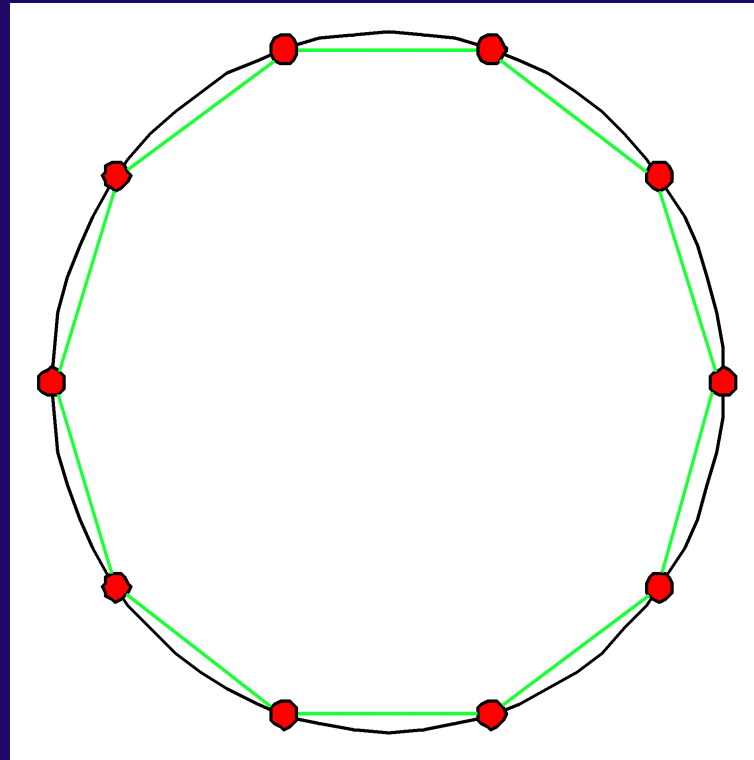
Jadranska 19, Ljubljana

Introduction

How to distribute n points equally on a circle?

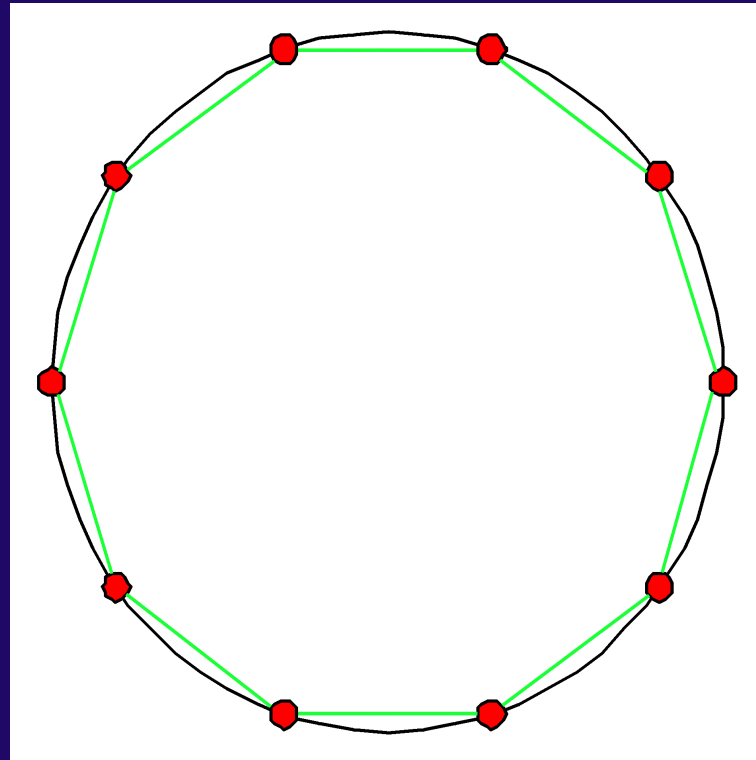
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But, how to distribute n points uniformly on a sphere?

Notation

We denote arrangement of n points with unit vectors

$$\mathbf{r}_1, \mathbf{r}_2, \dots, \mathbf{r}_n,$$

where

$$\mathbf{r}_i = (x_i, y_i, z_i)$$

is the position of the i -th point on the unit sphere S .

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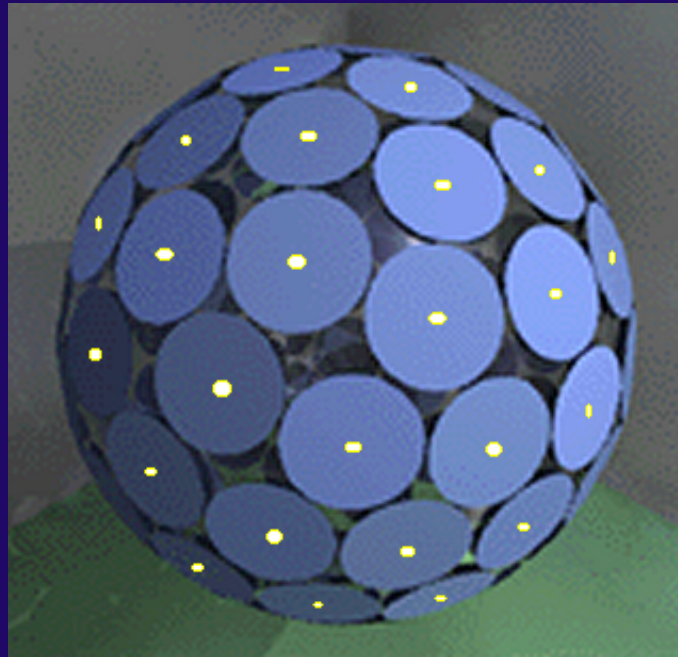
Equivalent to *the sphere packing problem* : find the largest diameter of n equal circles that can be placed on the sphere without overlap.

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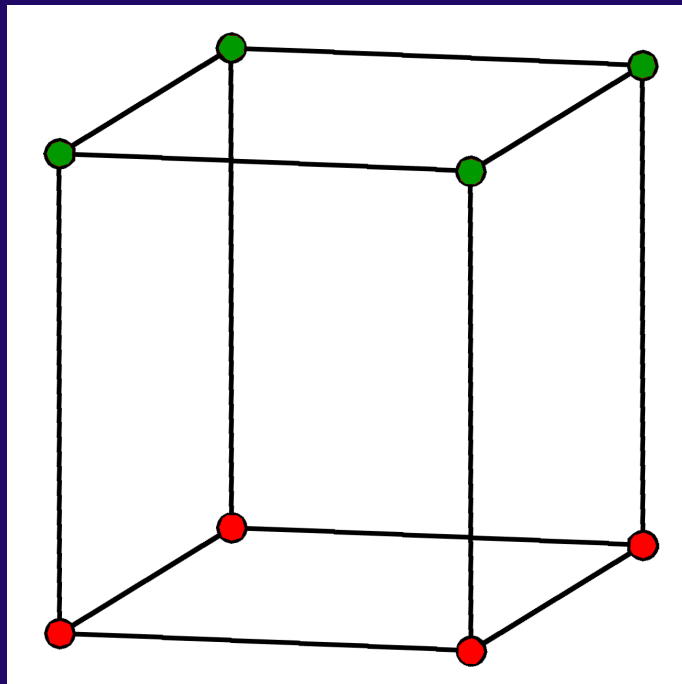
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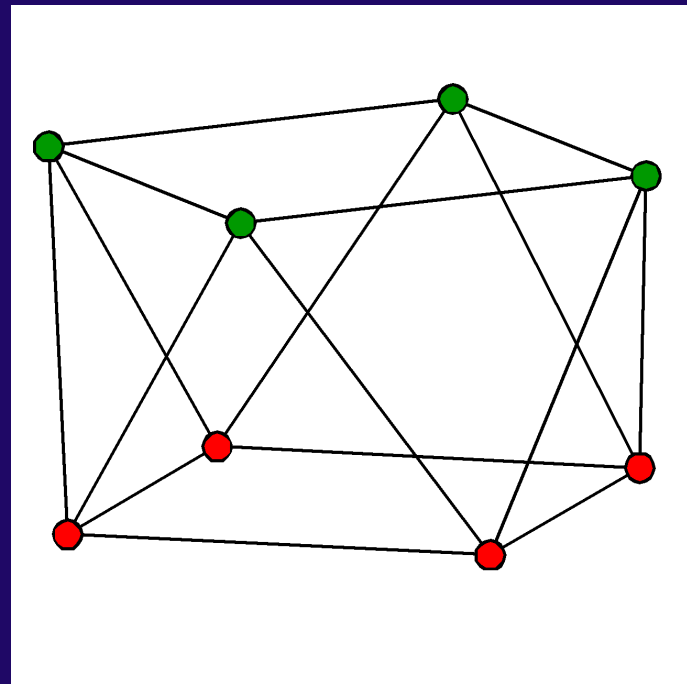
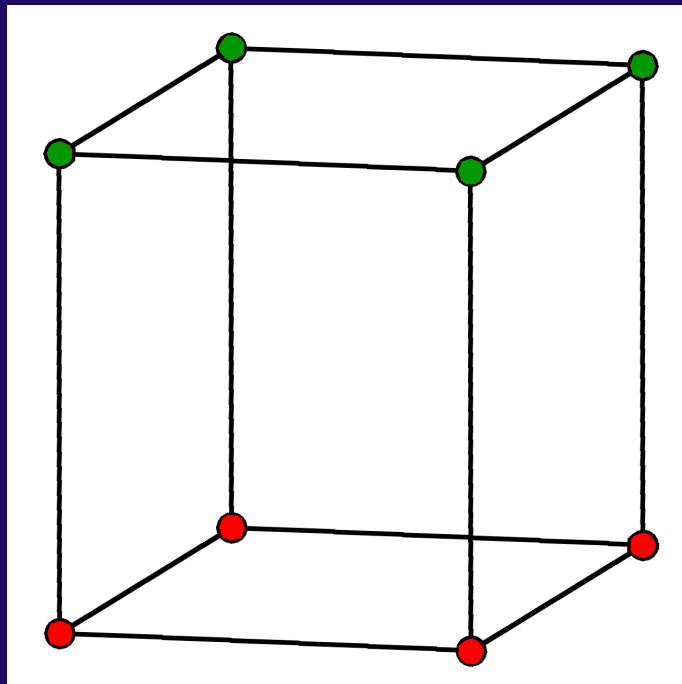


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- c) Use small random perturbations and search for the arrangement with greater minimum distance. When one is found, return to Step a).

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Find the minimum Coulomb potential

$$\sum_{i < j} \frac{1}{|\mathbf{r}_i - \mathbf{r}_j|}.$$

Also known as:

- *the electron problem,*
- *the Coulomb potential problem.*

The covering problem

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Energy function

For an arrangement \mathbf{R} and a given $\alpha \in \mathbb{R}$ we define α -energy

$$E(\mathbf{R}, \alpha) := \begin{cases} \sum_{1 \leq i < j \leq n} |\mathbf{r}_i - \mathbf{r}_j|^\alpha & \text{for } \alpha \neq 0 \\ \sum_{1 \leq i < j \leq n} \log \frac{1}{|\mathbf{r}_i - \mathbf{r}_j|} & \text{for } \alpha = 0. \end{cases}$$

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$\alpha = 0$: maximal product of distances between all pairs.

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Our computer program for Thompson's problem uses **gradient** method with the **random walk** method to escape from a local minimum. Two methods combine together in a kind of a **simulated annealing** method.

Known Tammes and Thompson exact solutions

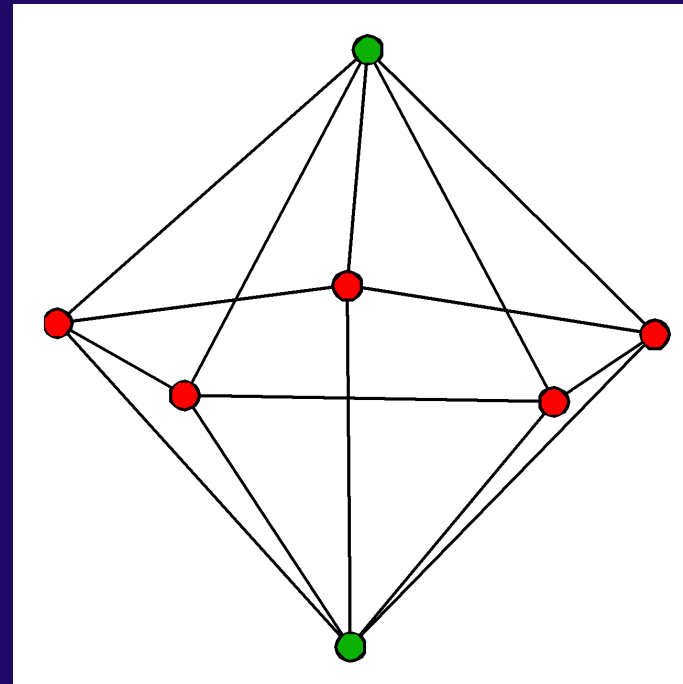
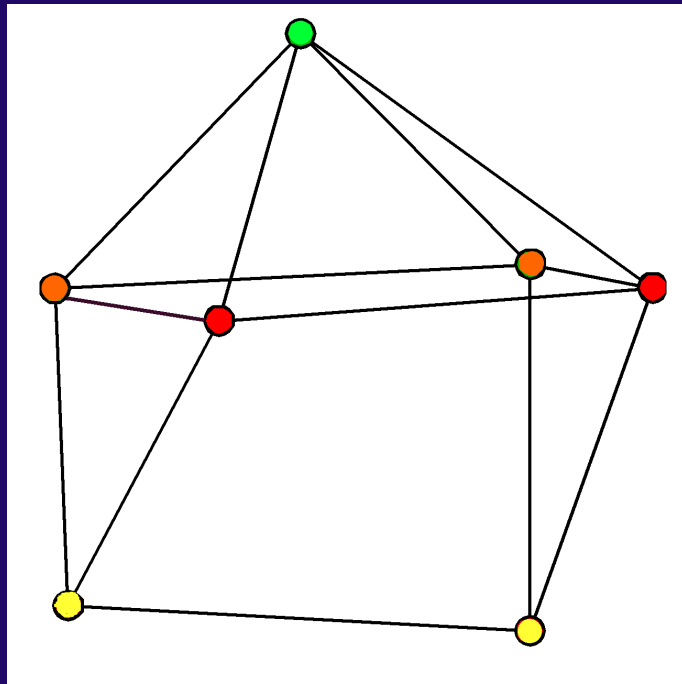
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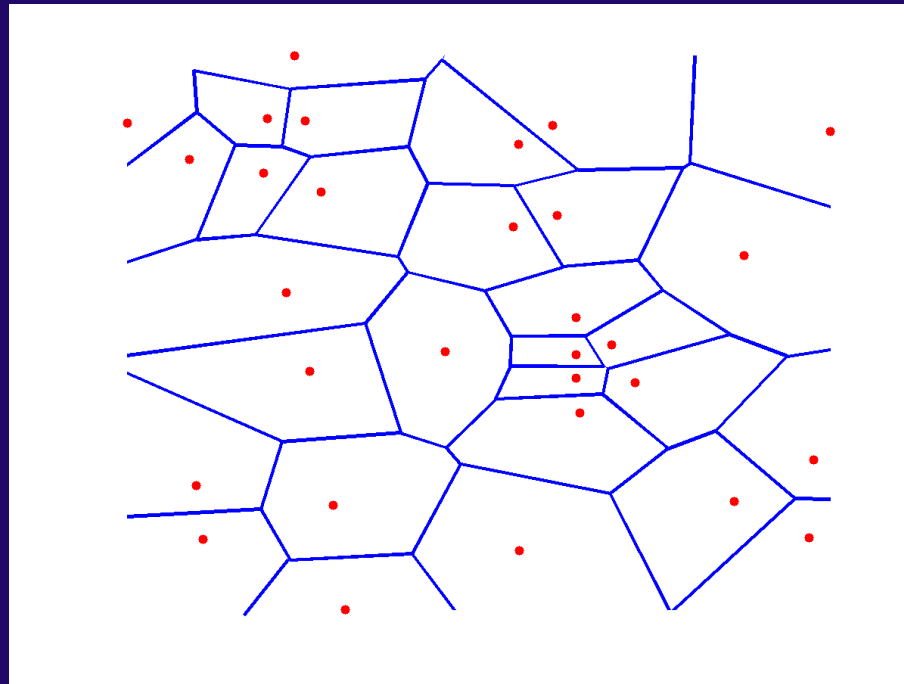
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Tammes and Thompson configuration for $n = 7$.

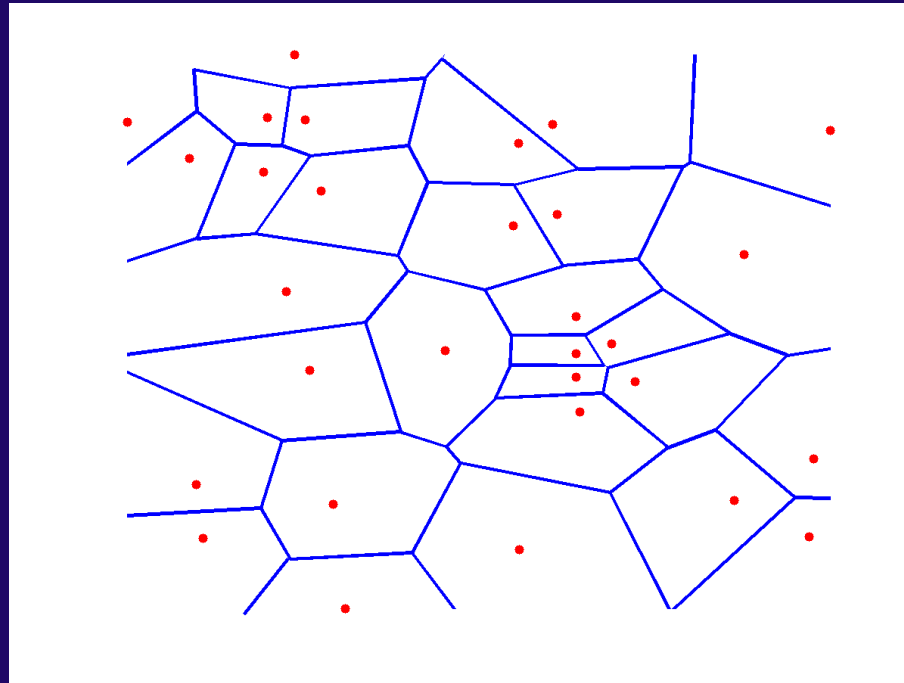
Voronoi polyhedra

To each arrangement \mathbf{R} we may associate a *Voronoi polyhedron*.

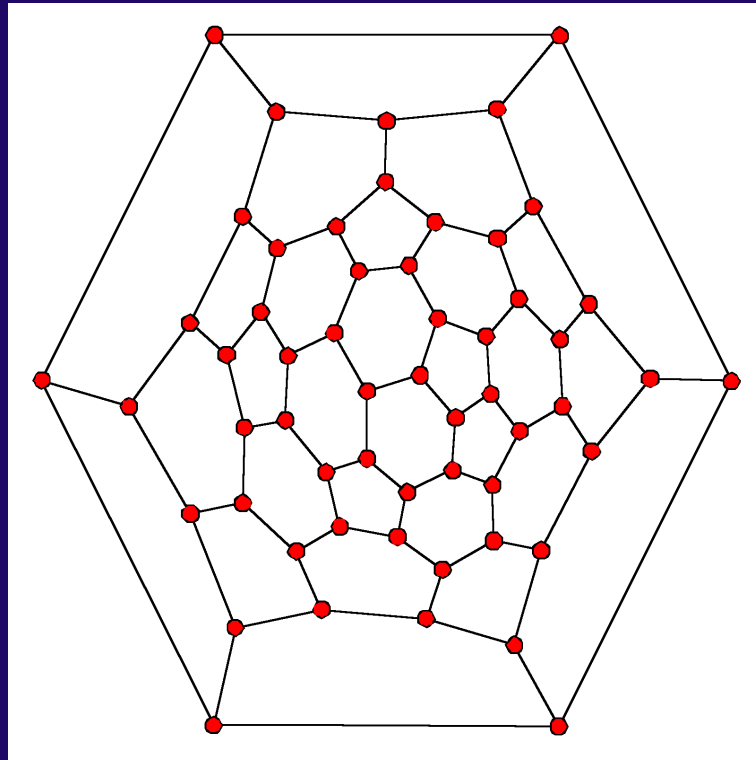


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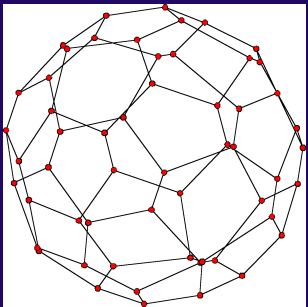
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For most n the associated Voronoi polyhedron is a fullerene.



The Schlegel diagram of Voronoi polyhedron of Thompson's arrangement for $n = 30$.



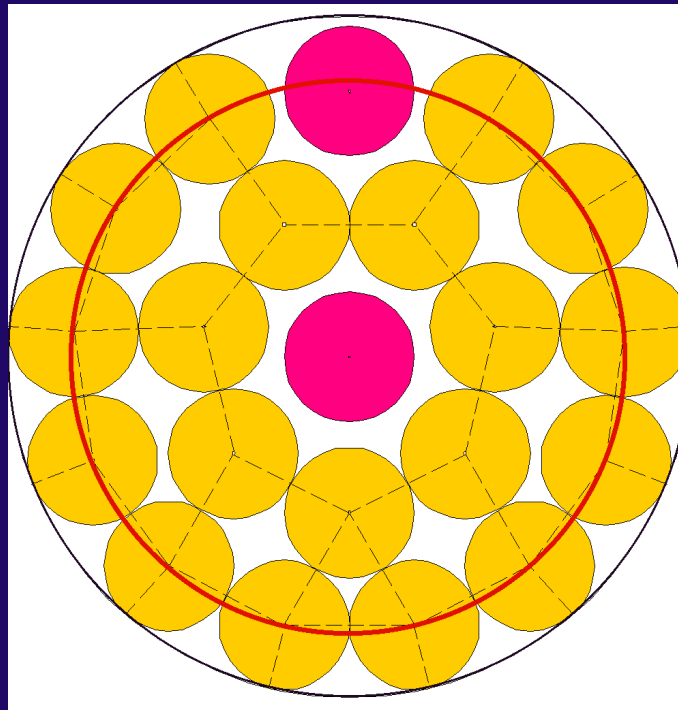
Uniform arrangements in a circle

Place n points inside the circle so that the minimal distance is as large as possible.

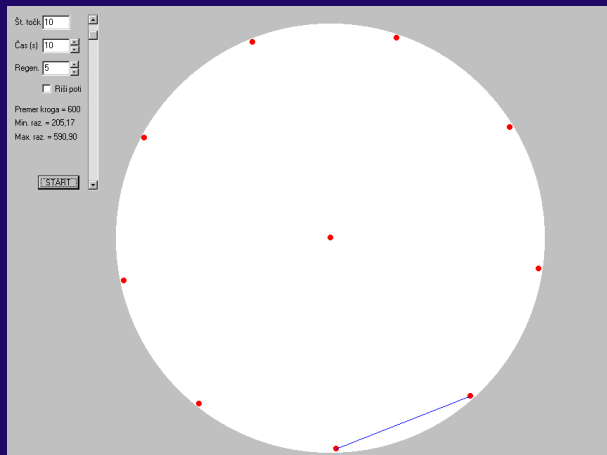
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This problem is equivalent to the problem of *packing of circles in a circle*.



The centers of circles represent the best arrangement of points in the circle going through the centers of the boundary circles.



by Peter Pogačar