

# Formulas for various domination numbers of products of paths and cycles

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# Polygraphs

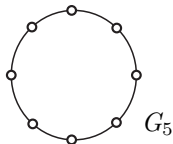
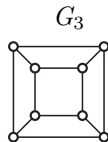
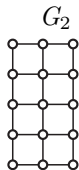
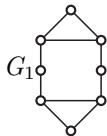
**Definition.** A *polygraph*  $\Omega_n = \Omega_n(G_1, \dots, G_n; X_1, \dots, X_n)$  over mutually disjoint monographs  $G_1, \dots, G_n$  has the vertex set

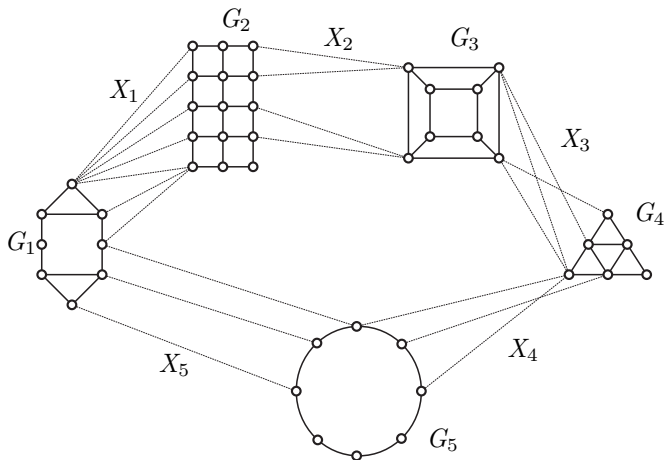
$$V(\Omega_n) = V(G_1) \cup \dots \cup V(G_n),$$

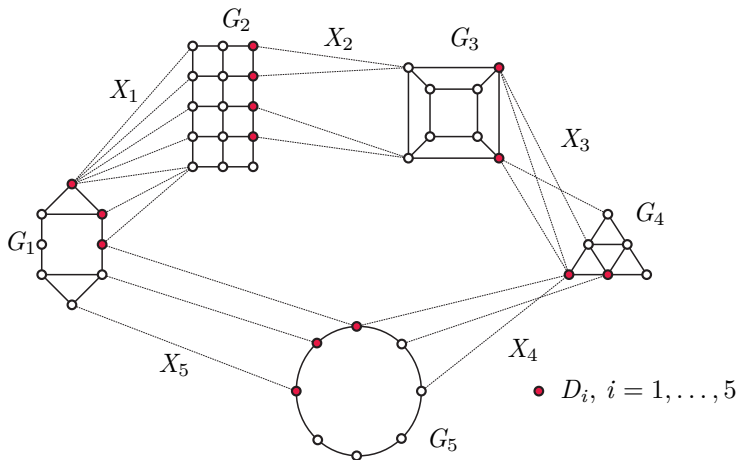
and the edge set

$$E(\Omega_n) = E(G_1) \cup X_1 \cup \dots \cup E(G_n) \cup X_n,$$

where  $X_i \subseteq V(G_i) \times V(G_{i+1})$  for  $i = 1, \dots, n$  and  $G_{n+1} \cong G_1$ .

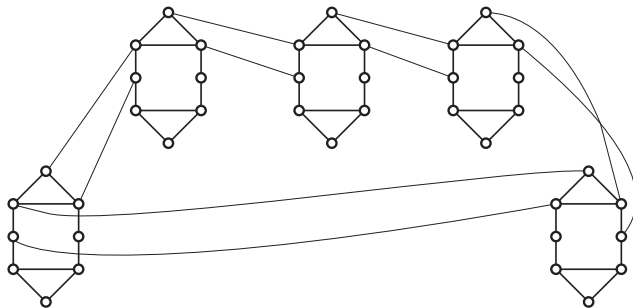




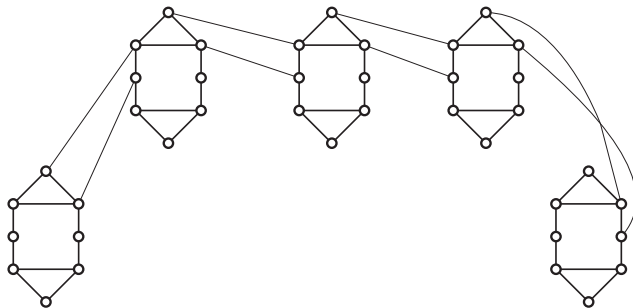




# Rotagraphs and fasciagraphs

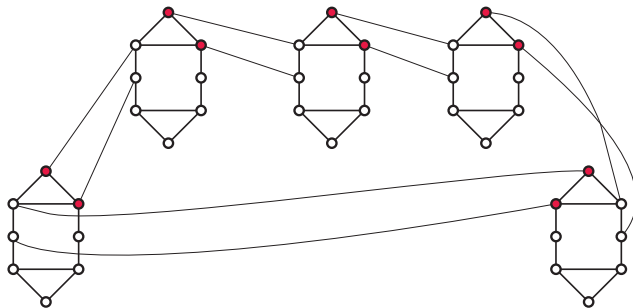


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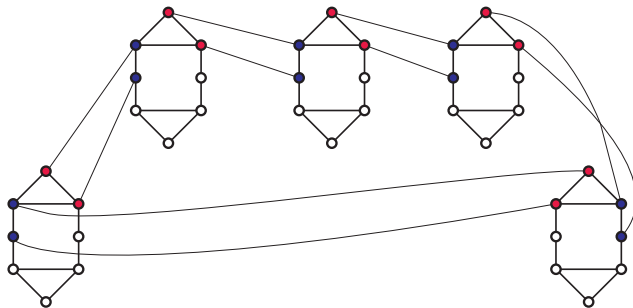


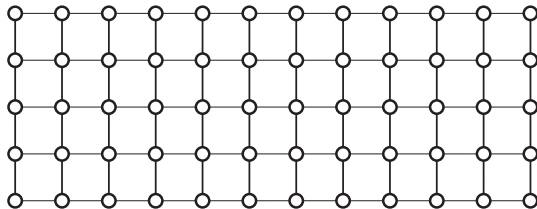


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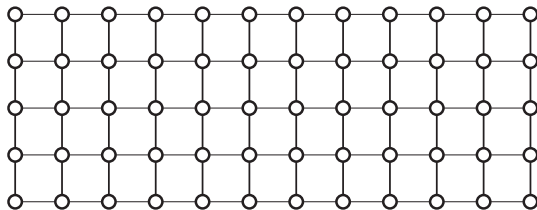
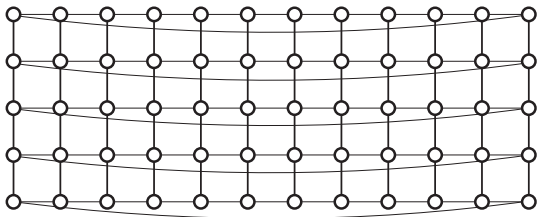


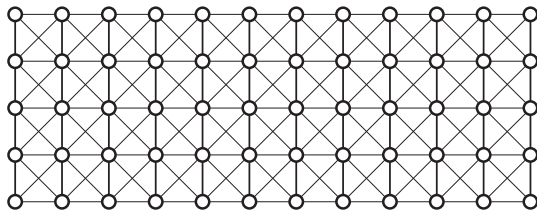
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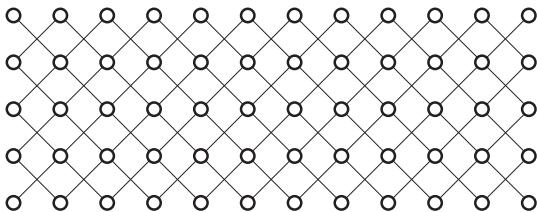


$$P_{12} \square P_5$$

 $P_{12} \square P_5$  $C_{12} \square P_5$



$$P_{12} \boxtimes P_5$$



$$P_{12} \times P_5$$

# The domination number and its variations

**Definition.** A set  $D \subseteq V$  of a graph  $G = (V, E)$  is a *dominating set*, if  $N[D] = V$ . The size of the smallest dominating set of a graph is the *domination number*,  $\gamma(G)$ .

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# Complexity results

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- bipartite graphs (Chang et al., 1984), chordal graphs (Booth and Johnson, 1985),...
- other domination types.

# Path algebra

**Definition.** A *semiring*  $\mathcal{P} = (P, \oplus, \circ, e^\oplus, e^\circ)$  is a set  $P$  on which two binary operations,  $\oplus$  and  $\circ$  are defined such that:

- 1  $(P, \oplus)$  is a commutative monoid with  $e^\oplus$  as a unit;
- 2  $(P, \circ)$  is a monoid with  $e^\circ$  as a unit;
- 3  $\circ$  is left- and right-distributive over  $\oplus$ ;
- 4 for every  $x \in P$ ,  $x \circ e^\oplus = e^\oplus = e^\oplus \circ x$

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## Examples of path algebras:

$$\mathcal{P}_1 = (\{0, 1\}, \max, \min, 0, 1)$$

$$\mathcal{P}_2 = (\mathbb{N}_0 \cup \{-\infty\}, \max, +, -\infty, 0)$$

$$\mathcal{P}_3 = (\mathbb{N}_0 \cup \{\infty\}, \min, +, \infty, 0)$$

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”tropical semiring”



## Matrices with elements of a path algebra:

Let  $\mathcal{P} = (P, \oplus, \circ, e^\oplus, e^\circ)$  be a path algebra and let  $\mathcal{M}_n(\mathcal{P})$  be the set of all  $n \times n$  matrices over  $P$ . Let  $A, B \in \mathcal{M}_n(\mathcal{P})$  and define operations  $\oplus$  and  $\circ$  in the usual way:

$$(A \oplus B)_{ij} = A_{ij} \oplus B_{ij},$$

$$(A \circ B)_{ij} = \bigoplus_{k=1}^n A_{ik} \circ B_{kj}.$$

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**Example:**  $\mathcal{P} = (\mathbb{N}_0 \cup \{\infty\}, \min, +, \infty, 0)$ .

- The zero matrix:

$$\begin{bmatrix} \infty & \infty & \dots & \infty \\ \vdots & \vdots & \ddots & \vdots \\ \infty & \infty & \dots & \infty \end{bmatrix}$$

- The unit matrix:

$$\begin{bmatrix} 0 & \infty & \dots & \infty \\ \infty & 0 & \dots & \infty \\ \vdots & \vdots & \ddots & \vdots \\ \infty & \infty & \dots & 0 \end{bmatrix}$$

- $(A \oplus B)_{ij} = \min \{A_{ij}, B_{ij}\}$ ,
- $(A \circ B)_{ij} = \min_{k=1, \dots, n} \{A_{ik} + B_{kj}\}$

# Path algebras and digraphs

Let  $\mathcal{P} = (P, \oplus, \circ, e^\oplus, e^\circ)$  be a path algebra and let  $G$  be a **labeled digraph**, that is a digraph together with a labeling function  $\ell$  which assigns to every arc of  $G$  an element of  $P$ :  $\ell : E(G) \rightarrow P$ . Let  $V(G) = \{v_1, v_2, \dots, v_n\}$  and define the following matrix:

$$(A(G))_{ij} = \begin{cases} \ell(v_i, v_j); & \text{if } (v_i, v_j) \text{ is an arc of } G \\ e^\oplus; & \text{otherwise} \end{cases}$$

The labeling  $\ell$  of  $G$  can be extended to **walks**: for a walk  $Q = (v_{i_0}, v_{i_1})(v_{i_1}, v_{i_2}) \dots (v_{i_{k-1}}, v_{i_k})$  of  $G$  let

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**Observation:** Let  $S_{ij}^k$  be the set of all walks of order  $k$  from  $v_i$  to  $v_j$  in  $G$ . Then

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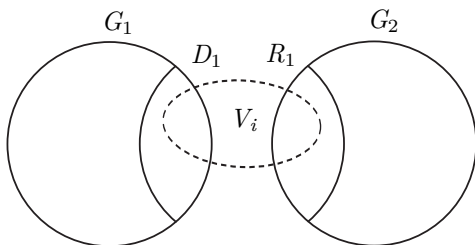
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# Algorithm, different graph invariants

$O(\log n)$  algorithm (Klavžar and Žerovnik, 1996)

Let  $\omega_n(G; X)$  be a rotagraph and  $\psi_n(G; X)$  a fasciagraph. Define a labeled digraph  $\mathcal{G} = \mathcal{G}(G; X)$ :

- $V(\mathcal{G}) \dots$  subsets of  $D \sqcup R$ ;
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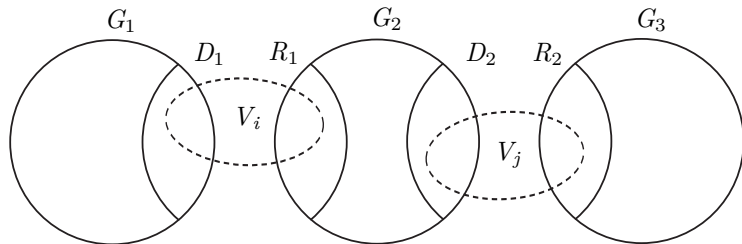
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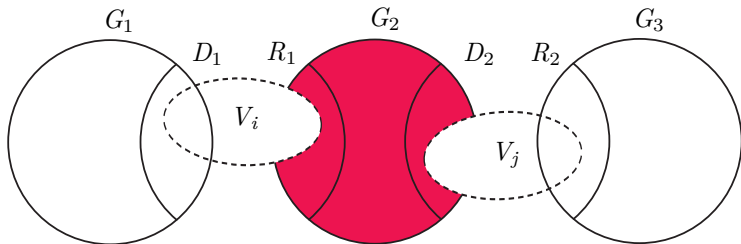
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- 3  $\ell(V_i, V_j) = |V_i \cap R| + \gamma_{i,j}(G; X) + |D \cap V_j| - |V_i \cap R \cap D \cap V_j|$ .



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Then

$$\gamma(\psi_n(G; X)) = (A(\mathcal{G})^n)_{00}$$

and

$$\gamma(\omega_n(G; X)) = \min_i (A(\mathcal{G})^n)_{ii}.$$

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- Time complexity;
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- Implementation to get some closed expressions.

## Calculating $A(\mathcal{G})^n$ in $O(1)$

Let  $A^n$  denote  $A(\mathcal{G})^n$  and  $A \in \mathcal{M}_n(\mathcal{P})$ , where  $\mathcal{P} = (\mathbb{N}_0 \cup \{\infty\}, \min, +, \infty, 0)$ .

### Lemma (-, Žerovnik, 2012)

Let  $N = |V(\mathcal{G}(G; X))|$ ,  $K = |V(G)|$ . Then there is an index  $q$  such that  $A^q = A^p + C$  for some index  $p < q$  and some constant matrix  $C = [c]_{ij}$ . Let  $P = q - p$ . Then for every  $r \geq p$  and every  $s \geq 0$  we have

$$A^{r+sP} = A^r + sC.$$

# Fasciagraphs

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$$A_{0i}^n = \min_k \{A_{0k}^{n-1} + A_{ki}\}.$$

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Assume that the  $j$ -th row of  $A^{n+P}$  and  $A^n$  differ for a constant,  $a_{ji}^{n+P} = a_{ji}^n + C$  for all  $i$ . Then  $\min_i a_{ji}^{n+P} = \min_i a_{ji}^n + C$ .

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$$\gamma(\psi_n(G; X)) = (A(\mathcal{G})^n)_{00}$$

$$A_{0i}^n = \min_k \{A_{0k}^{n-1} + A_{ki}\}.$$

## Lemma (-, Žerovnik, 2012)

Assume that the  $j$ -th row of  $A^{n+P}$  and  $A^n$  differ for a constant,  $a_{ji}^{n+P} = a_{ji}^n + C$  for all  $i$ . Then  $\min_i a_{ji}^{n+P} = \min_i a_{ji}^n + C$ .

# Rotagraphs

## Lemma (-, Žerovnik, 2012)

Let  $A^q = A^p + C$  and  $P = q - p$ . Then for every  $t \in 0, 1, \dots, P - 1$  there is a constant  $C_t$  such that for all  $n \geq p$  with  $t \equiv (n - p) \pmod{P}$  we have

$$\gamma(\psi_n(G; X)) - \gamma(\omega_n(G; X)) = C_t.$$

# Results and remarks

## Theorem (-, Žerovnik, 2012)

Domination numbers of fasciagraphs and rotagraphs can be computed in constant time, i.e. independently of the size of a monograph  $G$ .



$k$	$\gamma(P_n \square C_k)$
<b>3</b>	$\left\lceil \frac{3n}{4} \right\rceil + 1;$ <b>if</b> $n \equiv 0 \pmod{4}$ $\left\lceil \frac{3n}{4} \right\rceil;$ <b>otherwise</b>
<b>4</b>	$n$
<b>5</b>	<b>3;</b> <b>if</b> $n = 2$ <b>4;</b> <b>if</b> $n = 3$ $n + 2;$ <b>otherwise</b>
<b>6</b>	$\left\lceil \frac{4n}{3} \right\rceil;$ <b>if</b> $n \equiv 1 \pmod{3}$ $\left\lceil \frac{4n}{3} \right\rceil + 1;$ <b>otherwise</b>
<b>7</b>	$\left\lceil \frac{3n}{2} \right\rceil + 1;$ <b>if</b> $n \equiv 1 \pmod{2}$ $\left\lceil \frac{3n}{2} \right\rceil + 2;$ <b>otherwise</b>

$k$	$\gamma(P_n \square C_k)$
<b>8</b>	<b>4;</b> <b>if</b> $n = 2$
	<b>6;</b> <b>if</b> $n = 3$
	<b>8;</b> <b>if</b> $n = 4$
	$\left\lfloor \frac{9n}{5} \right\rfloor + 1;$ <b>if</b> $n \equiv 5 \pmod{10}$
	$\left\lfloor \frac{9n}{5} \right\rfloor + 2;$ <b>otherwise</b>
<b>9</b>	<b>5;</b> <b>if</b> $n = 2$
	<b>7;</b> <b>if</b> $n = 3$
	<b>10</b> <b>if</b> $n = 4$
	$2n + 2;$ <b>otherwise</b>
<b>10</b>	$2n + 2;$ <b>if</b> $n \leq 5$
	$2n + 3;$ <b>if</b> $6 \leq n \leq 9$
	$2n + 4$ <b>otherwise</b>
<b>11</b>	$\left\lfloor \frac{19n}{8} \right\rfloor + 1;$ <b>if</b> $n \in \{1, 2, 4, 6\}$ <b>or</b> $n \equiv 3 \pmod{8}$
	$\left\lfloor \frac{19n}{8} \right\rfloor + 2;$ <b>otherwise</b>

$k$	$\gamma(C_n \square P_k)$	$\gamma(C_n \square C_k)$
2	$\begin{cases} \left\lceil \frac{n}{2} \right\rceil + 1; & \text{if } n \equiv 2 \pmod{4} \\ \left\lceil \frac{n}{2} \right\rceil; & \text{otherwise} \end{cases}$	
3	$\left\lceil \frac{3n}{4} \right\rceil$	$\left\lceil \frac{3n}{4} \right\rceil$
4	$\begin{cases} n + 1; & \text{if } n \in \{5, 9\} \\ n; & \text{otherwise} \end{cases}$	$n$
5	$\begin{cases} 4; & \text{if } n = 3 \\ \left\lceil \frac{6n}{5} \right\rceil + 1; & \text{if } n \equiv 3, 5, 9 \pmod{10} \\ \left\lceil \frac{6n}{5} \right\rceil; & \text{otherwise} \end{cases}$	$\begin{cases} n; & \text{if } n \equiv 0 \pmod{5} \\ n + 2; & \text{if } n \equiv 3 \pmod{5} \\ n + 1; & \text{otherwise} \end{cases}$
6	$\begin{cases} 9; & \text{if } n = 6 \\ \left\lceil \frac{10n}{7} \right\rceil + 1; & n \equiv 2, 6, 7, 9 \\ & 13 \pmod{14} \\ \left\lceil \frac{10n}{7} \right\rceil; & \text{otherwise} \end{cases}$	$\begin{cases} \left\lceil \frac{4n}{3} \right\rceil + 1; & n \equiv 2, 3, 8, 9 \pmod{18} \\ & 11, 14, 15, 17 \pmod{18} \\ \left\lceil \frac{4n}{3} \right\rceil; & \text{otherwise} \end{cases}$
7	$\begin{cases} 6; & \text{if } n = 3 \\ 16; & \text{if } n = 9 \\ 36; & \text{if } n = 21 \\ \left\lceil \frac{5n}{3} \right\rceil; & \text{otherwise} \end{cases}$	$\begin{cases} \left\lceil \frac{3n}{2} \right\rceil; & n \equiv 0, 5, 9 \pmod{14} \\ \left\lceil \frac{3n}{2} \right\rceil + 2; & n \equiv 2, 8, 12 \pmod{14} \\ \left\lceil \frac{3n}{2} \right\rceil + 1; & \text{otherwise} \end{cases}$